

# Moment and Probability Inequalities

Statistics 110

Summer 2006



# Moment Inequalities

- Schwarz's Inequality (sometimes called Cauchy-Schwarz)

$$(E[XY])^2 \leq E[X^2]E[Y^2]$$

**Proof.** Suppose that  $E[X^2] > 0$  and  $E[Y^2] > 0$  Let

$$U = \frac{X}{\sqrt{E[X^2]}} \quad \text{and} \quad V = \frac{Y}{\sqrt{E[Y^2]}}$$

It can be shown that  $2|UV| \leq U^2 + V^2$ . Thus

$$2|E[UV]| \leq 2E[|UV|] \leq E[U^2] + E[V^2] = 2$$

This gives

$$(E[UV])^2 \leq (E[|UV|])^2 \leq 1$$

implying

$$\frac{(E[XY])^2}{E[X^2]E[Y^2]} \leq \frac{(E[|XY|])^2}{E[X^2]E[Y^2]}$$

□

One consequence of this inequality is that  $(\text{Cov}(X, Y))^2 \leq \text{Var}(X)\text{Var}(Y)$  or  $|\text{Cov}(X, Y)| \leq \sigma_X\sigma_Y$ . A consequence of this is that  $|\text{Corr}(X, Y)| \leq 1$ , a result discussed earlier.

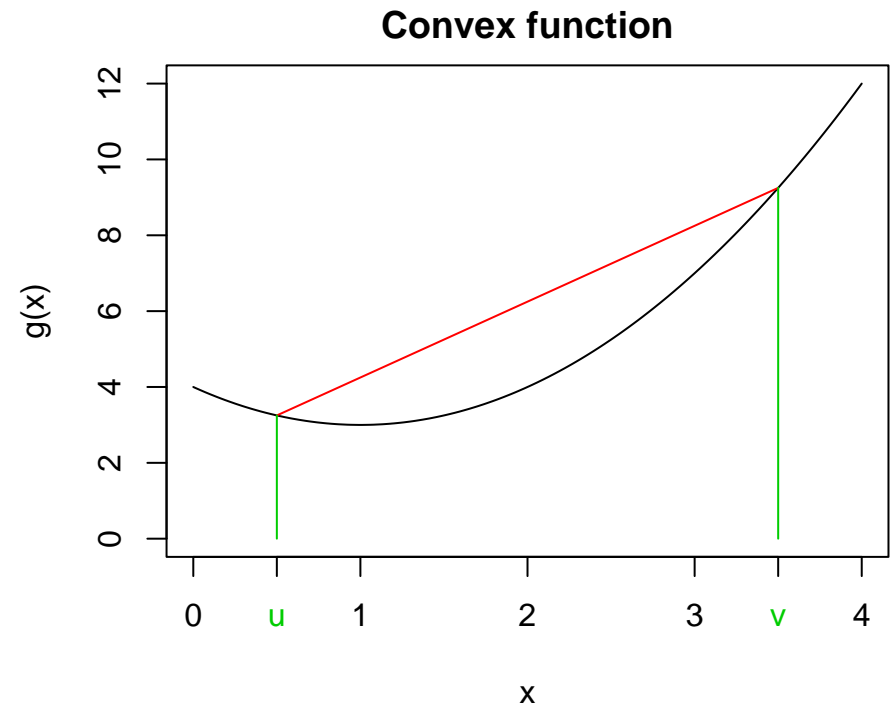
- Jensen's Inequality

If  $g(\cdot)$  is a convex function on the interval  $(a, b)$  and  $X$  is a RV taking values in  $(a, b)$ , then  $E[g(X)] \geq g(E[X])$ .

Note that a function  $g(\cdot)$  is convex on the open interval  $I = (a, b)$  if

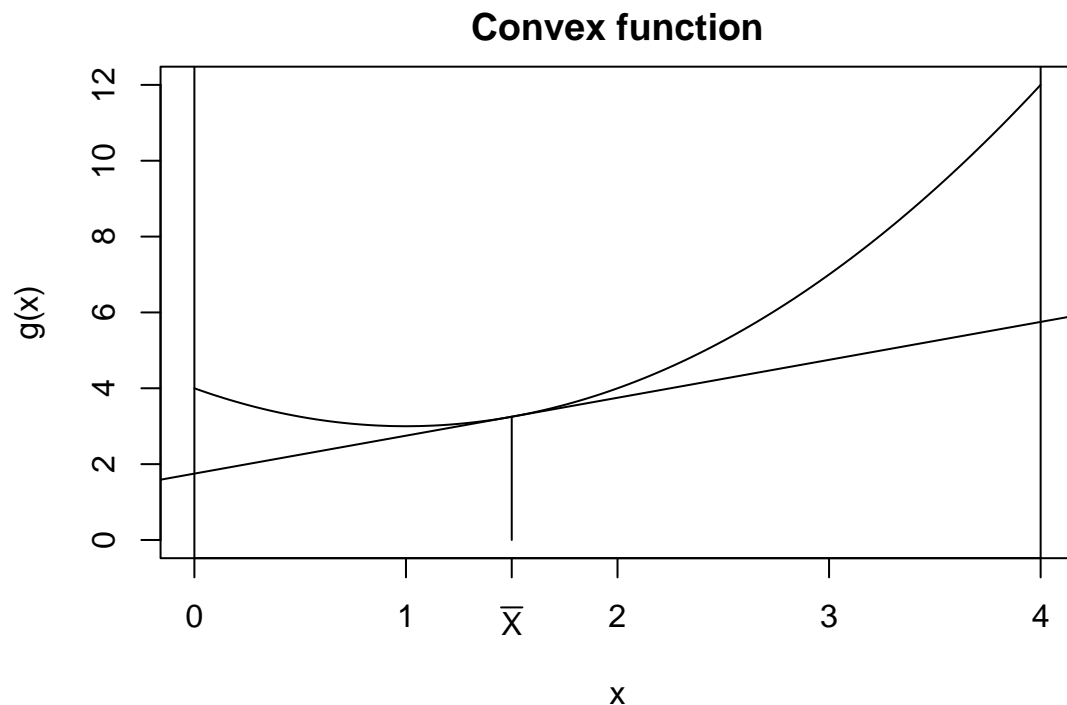
$$g(\alpha u + (1 - \alpha)v) \leq \alpha g(u) + (1 - \alpha)g(v)$$

for all  $u, v \in I$  and  $0 \leq \alpha \leq 1$ .



**Proof.** Convexity means that a supporting line exists at each  $t \in (a, b)$ . i.e. the graph lies completely above each tangent line.

From the supporting line at  $t = E[X]$  (with slope  $\lambda$ ), we have



$$\begin{aligned}
 g(x) &\geq g(E[X]) + \lambda(x - E[X]) \\
 E[g(X)] &\geq E[g(E[X]) + \lambda(X - E[X])] \\
 &= g(E[X]) + \lambda(E[X] - E[X]) = g(E[X])
 \end{aligned}$$

□

A couple of examples where Jensen's inequality can be used are the following

1.  $E[e^X] \geq \exp(E[X]).$

For example, assume  $X \sim N(\mu, \sigma^2)$  and let  $Y = e^X \sim \text{logN}(\mu, \sigma^2).$

A consequence is that  $E[Y] = E[e^X] \geq e^\mu.$

In fact  $E[Y] = e^{\mu+0.5\sigma^2}$

Note going the other way, we get  $\log(E[X]) \geq E[\log X]$  since  $-\log x$  is a convex function ( $\log x$  is a concave function).

i.e.  $\log e^{\mu+0.5\sigma^2} \geq \mu$

## 2. Arithmetic mean $\geq$ Geometric Mean $\geq$ Harmonic Mean

For any set of  $n$  positive numbers  $x_1, x_2, \dots, x_n$ ,

$$\frac{x_1 + \dots + x_n}{n} \geq \sqrt[n]{x_1 x_2 \dots x_n} \geq \frac{n}{\frac{1}{x_1} + \dots + \frac{1}{x_n}}$$

To justify the first inequality let  $X$  be a random variable taking values  $x_1, x_2, \dots, x_n$  each with probability  $\frac{1}{n}$ . Then Jensen's says

$$\log \left( \frac{x_1 + \dots + x_n}{n} \right) \geq \frac{\log x_1 + \dots + \log x_n}{n} = \log(x_1 \dots x_n)^{1/n}$$

Then exponentiate both sides to get the first inequality.

The other inequalities can be derived similarly.

- Lyapunov's Inequality

If  $0 < s < t$

$$(E[|X|^s])^{1/s} \leq (E[|X|^t])^{1/t}$$

A consequence of this is the relationship (for some integer  $p$ )

$$E[|X|] \leq (E[|X|^2])^{1/2} \leq (E[|X|^3])^{1/3} \leq \dots \leq (E[|X|^p])^{1/p}$$

which implies

$$|E[X]|^q \leq (E[|X|])^q \leq E[|X|^q] \text{ if } 1 \leq q \leq p$$

**Proof.** Let  $r = \frac{t}{s} > 1$ . Let  $Y = |X|^s$  and apply Jensen's inequality to  $g(y) = |y|^r$ , giving  $(E[|Y|])^r \leq E[|Y|^r]$ . This implies that

$$(E[|X|^s])^{t/s} \leq E[|X|^t]$$

Taking the  $t$ th root of each side gives the result.  $\square$



# Probability Inequalities

- Markov Inequality

Let  $X$  be a non-negative RV (i.e.  $P[X \geq 0] = 1$ ). Then for any  $a > 0$ ,

$$P[X \geq a] \leq \frac{E[X]}{a}$$

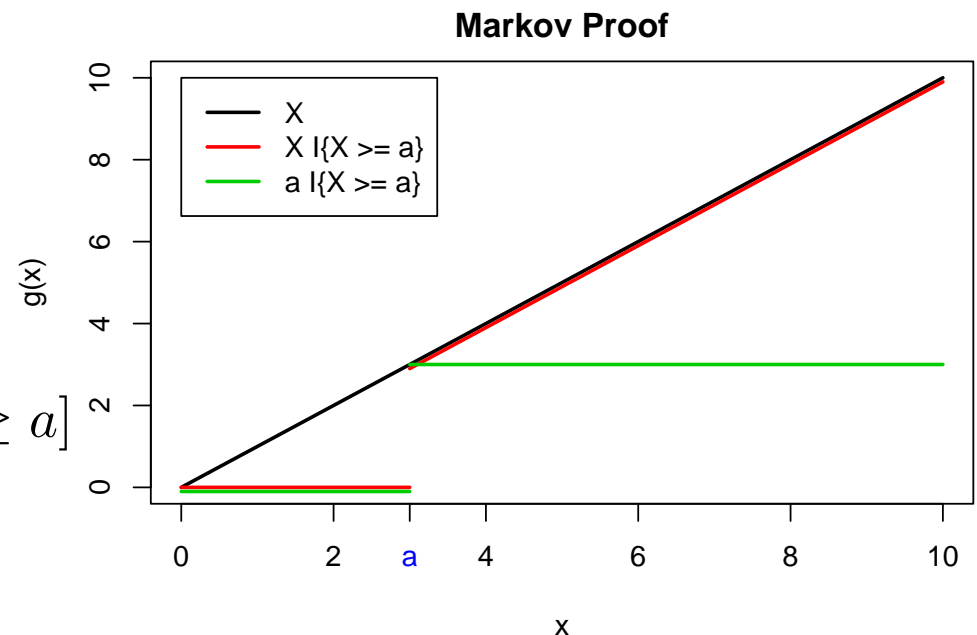
## Proof.

$$X \geq XI\{X \geq a\} \geq aI\{X \geq a\}$$

Therefore

$$E[X] \geq E[XI\{X \geq a\}] = aP[X \geq a]$$

□



Note that there is an alternative version of this inequality that says if  $E[X^r] < \infty$ ,

$$P[X \geq a] \leq \frac{E[X^r]}{a^r}$$

- Chebyshev's Inequality.

If  $E[X] = \mu$  and  $\text{Var}(X) = \sigma^2 < \infty$ , then

$$P[|X - \mu| \geq k] \leq \frac{\sigma^2}{k^2}$$

Note that this equality is sometimes written as the equivalent

$$P[|X - \mu| \geq k\sigma] \leq \frac{1}{k^2}$$

## Proof.

$$P[|X - \mu| \geq a] = P[(X - \mu)^2 \geq a^2] \leq \frac{E[(X - \mu)^2]}{a^2} \quad \text{by Markov's Inequality}$$

Take  $a = k$  to get the first form of the result and  $a = k\sigma$  to get the second form of the result.  $\square$

Example: Suppose it is known that the number of widgets produced for Guinness breweries in a factory during an hour is a RV with mean 500.

1. What can be said about the probability that an hour's production will exceed 1000?

Answer: By Markov's inequality

$$P[X \geq 1000] \leq \frac{E[X]}{1000} = \frac{500}{1000} = 0.5$$

2. If the variance of a hour's production is known to be 100, then what can be said about the probability that a hour's production will be between 450 and 550?

Answer: By Chebyshev's inequality

$$P[|X - 500| \geq 50] \leq \frac{\text{Var}(X)}{50^2} = \frac{100}{50^2} = \frac{1}{25} = 0.04$$

This implies that

$$P[|X - 500| < 50] \geq 1 - \frac{1}{25} = \frac{24}{25} = 0.96$$

3. What can be said about the probability that the production will be between 450 and 550 if  $X$  is normally distributed ( $N(500, 100)$ )?

$$\begin{aligned} P[450 \leq X \leq 550] &= P\left[\frac{450 - 500}{10} \leq Z \leq \frac{550 - 500}{10}\right] \\ &= P[-5 \leq Z \leq 5] = \Phi(5) - \Phi(-5) = 0.9999994 \end{aligned}$$

Note that these bounds are not particularly tight in most cases.

In fact they are what happens in a “worst case scenario”.

The following inequality also fits into this setting, where the bounds are often loose.

- One-sided Chebyshev's Inequality

If  $E[X] = \mu$  and  $\text{Var}(X) = \sigma^2 < \infty$ , then for any  $a > 0$ ,

$$P[X \geq \mu + a] \leq \frac{\sigma^2}{\sigma^2 + a^2}$$

$$P[X \leq \mu - a] \leq \frac{\sigma^2}{\sigma^2 + a^2}$$

**Proof.** Without loss of generality, assume that  $\mu = 0$ . Then for any  $b$ ,

$$\begin{aligned} P[X \geq a] &= P[X + b \geq a + b] \\ &= P[(X + b)^2 \geq (a + b)^2] \\ &\leq \frac{E[(X + b)^2]}{(a + b)^2} = \frac{E[X^2] + b^2}{(a + b)^2} \\ &= \frac{\alpha + t^2}{(1 + t)^2} \stackrel{\text{Def}}{=} g(t) \end{aligned}$$

where

$$\alpha = \frac{E[X^2]}{a^2} = \frac{\sigma^2}{a^2}; \quad t = \frac{b}{a}$$

To minimize  $g(t)$  (i.e. find the best  $b$ ), set  $t = \alpha$ , yielding

$$\min g(t) = \frac{\alpha + \alpha^2}{(1 + \alpha)^2} = \frac{\frac{\sigma^2}{a^2}}{1 + \frac{\sigma^2}{a^2}} = \frac{\sigma^2}{\sigma^2 + a^2}$$

The other inequality is proved similarly.

□

Example: Back to the widget example. What can be said about the probability that last least 550 widgets are made, assuming the mean is 500 and the variance is 100?

Answer:

$$P[X \geq 550] = P[X \geq 500 + 50] \leq \frac{\sigma^2}{\sigma^2 + 50^2} = \frac{100}{100 + 2500} = 0.0384$$

If we only use the different forms of the Markov inequality we get

$$P[X \geq 550] \leq \frac{E[X]}{550} = \frac{500}{550} = 0.909$$

and

$$P[X \geq 550] \leq \frac{E[X^2]}{550^2} = \frac{\sigma^2 + \mu^2}{550^2} = 0.827$$

Note that if the production was normally distributed,  $P[X \geq 550] = 0.000000287$

These probability bounds may not be useful as they may give values greater than 1. For example, if  $\mu = 500$ , the Markov bound for

$$P[X \geq 400] \leq \frac{500}{400} = 1.25$$

This is a reason why different bounds have been developed. Generally, the stronger the assumptions you make, the tighter the bounds you can get.